

1 A multi-product and multi-period supply chain network design
2 problem with price-sensitive demand and incremental quantity
3 discount

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10 **Abstract**

11 Demand is an important factor in supply chain network design (SCND). Among the many factors
12 that affect demand, the impact of price cannot be underestimated. This study, for the first time,
13 provides enterprises with direct multi-period pricing and demand decision solutions in SCND by si-
14 multaneously considering the impact of price-sensitive demand and incremental quantity discounts
15 on the network. In this way, a mixed-integer nonlinear programming model is established based on
16 the relationships between price and demand functions to maximize the overall profit of the supply
17 chain. Due to the nonlinear characteristics of the model, this study proves the existence of the
18 optimal solution of the problem and the feasibility of implementing the following two algorithms by
19 analyzing the nature of the problem. The first algorithm is based on the second-order cone program-
20 ming (SOCP) method, and the second algorithm is based on the outer-approximation method. The
21 results of numerical experiments at different scales show that the SOCP method can obtain the opti-
22 mal solution for small and medium-scale experiments. In contrast, the outer-approximation method
23 can find approximate optimal solutions within a reasonable time for large-scale experiments. The
24 results confirm that considering incremental quantity discount can significantly enhance the prof-
25 itability of supply chain networks in long-term planning. Finally, some future research directions in
26 the field of SCND are discussed.

27 *Keyword:* supply chain network design; price-sensitive demand; incremental quantity discount;
28 outer-approximation; second-order cone programming

29 **1. Introduction**

30 In the supply chain network design (SCND) problems, a multitude of factors must be taken into con-
31 sideration, including the selection of factory facilities, establishment of warehouses, transportation
32 of materials, and more. However, within this array of factors, the pivotal role of customer demand
33 cannot be overstated (Ardjmand et al., 2016). This is because demand directly restricts sales vol-
34 ume, which in turn controls the quantities of goods produced and transported. In the context of
35 a demand-driven market operating model, customer demand dictates the entirety of material flow
36 within the supply chain network, as well as the corresponding facility siting. Therefore, demand
37 drives production activities and imposes an essential influence on the profit of the entire supply
38 chain network.

39 While numerous scholars acknowledge the paramount role of demand in SCND, in traditional
40 SCND problems, demand is often treated as a given parameter. In reality, demand is influenced by
41 various factors, including market dynamics, governmental policies, and environmental conditions.
42 Among these factors, the demand for the majority of products is contingent upon price decisions
43 (Hanjoul et al., 1990), thus constituting price-sensitive demand. Although the profit for per unit of
44 each product increases when the price rises, the customer demand decreases, and consequently the

45 total profit does not necessarily improve. An example is the food supply chain. According to the
46 report released by the Food and Agriculture Organization of the United Nations in November 2022,
47 the palm oil price was affected by the COVID-19 pandemic and extreme weather in recent years.
48 As a result, the top palm oil consumers in India market are shifting to purchase relatively lower
49 priced soybean oil and even sunflower oil. When the demand depends on the price of the commodity,
50 inventory and selling price should be jointly decided, and finding the optimal pricing strategy is one
51 of the main strategies to obtain the maximum profit (Ghoreishi et al., 2015). Therefore, it is of both
52 theoretical and practical significance to consider the impact of price decisions on consumer demands
53 in supply chain network research.

54 In the integrated optimization of the supply chain network, it avoids the independent decision-
55 making of the supply chain members and only optimizes their own goals, but develops a coordination
56 mechanism to coordinate the decisions of the chain members to achieve high or optimal chain per-
57 formance. Among various coordination mechanisms, quantity discounts have been proven to be
58 successful coordination mechanisms within supply chains (Hanh et al., 2022). Discount policies of-
59 fered by suppliers can harmonize the supply chain and generate higher or even maximum supply
60 chain profits in a price sensitive environment (Hsieh et al., 2010). Quantity discounts are an effective
61 mechanism for coordinating decentralized supply chains, reducing inventory-related costs within the
62 supply chain and potentially attracting more demand (Ke & Bookbinder, 2018). From the buyer's
63 point of view, quantity discounts can calculate the optimal order quantity that minimizes the buyer's
64 total associated costs. However, switching to the supplier's point of view, quantity discounts may
65 be used to attract buyers to increase the order quantity and improve the supplier's profit.

66 In recent years, many scholars have also paid attention to the impact of price-sensitive demand
67 and quantity discount on SCND (Ahmadi-Javid & Ghandali, 2014; Fattahi et al., 2015; Ghoreishi
68 et al., 2015; Huang et al., 2018). However, there is still room for further exploration in the current
69 research in terms of practicality. The first is that price and demand are not treated simultaneously
70 as variables in network design. Ahmadi-Javid & Hoseinpour (2015a) studied for the first time the
71 profit maximization problem in the location selection inventory problem considering price-sensitive
72 demand, but only using price as a decision variable. Although Mogale et al. (2022) regarded price
73 and demand as decision variables at the same time, its price is not a multi-period dynamic decision,
74 and quantity discount is not considered. Second, there are fewer studies considering both price-
75 sensitive demand and quantity discount, especially incremental quantity discount. Previous studies
76 have mostly used price and demand as parameters, or only one of them as variables. However, in
77 the price-sensitive demand relationship, the uncertainty of the two can truly reflect the needs of
78 customers in different pricing. Therefore, our paper selects product price and demand as decision
79 variables, while considering the impact of incremental quantity discount on the entire supply chain.
80 This approach provides a more instructive decision plan for SCND problems.

81 This paper studies an SCND with price-sensitive demand, and establishes a three-echelon,
82 multi-product, and multi-period supply chain network structure consisting of multiple manufactur-
83 ers, multiple warehouses, and a single customer group. The impact of incremental quantity discount
84 on network profitability is considered at the suppliers. In addition to the decision variables (location
85 and production decisions as well as inventory and transportation decisions) that are considered in
86 traditional SCND problems, this study also classifies prices as decision variables and establishes a
87 mixed-integer nonlinear programming (MINLP) model with the objective to maximize the profit
88 of the supply chain network. The model takes into account constraints such as production, trans-
89 portation, and inventory capacity of the companies in the supply chain, as well as the flow balance
90 between production and transportation. Furthermore, the constraints of price-demand relationships
91 are added. Consequently, the established model can more accurately reflect changes of the de-
92 mand in response to different price decisions. Based on this, we would like to answer the following
93 questions:

- 94 (1) Which manufacturers and warehouses should be open to producing and storing products in
95 SCND problems?
- 96 (2) How much is the volume of finished goods produced, stored, and transported between network-
97 level facilities? Which discount should upstream manufacturers choose when purchasing raw
98 materials and how much raw materials are purchased at different discount levels?
- 99 (3) What is the impact of prices and demand for different products and incremental quantity dis-
100 count on network profits in different periods?

101 The remainder of the paper is structured as follows: Section 2 reviews the literature related to
102 this study. Section 3 describes the incremental quantity discounts used in this paper and constructs
103 an SCND model that takes into account price-sensitive demand. Section 4 uses the SOCP and the
104 outer-approximation method to address the nonlinearity of the model. Section 5 verifies the validity
105 and feasibility of the algorithms by computational experiments. This section also analyzes the model
106 without considering discounts to further specify the advantages of the model proposed in this paper.
107 Finally, Section 6 presents relevant conclusions and future research directions.

108 **2. Literature review**

109 Problems of traditional SCND have been studied in depth over the past few decades. Several studies,
110 such as Peng et al. (2011), De Matta & Miller (2015), and Umpfenbach et al. (2018), focused on
111 strategic planning decisions that include capacity and supplier selection decisions of SCND. Geor-
112 giadis et al. (2011), Ardjmand et al. (2016), Behzadi et al. (2017), Li et al. (2017), Pervin et al.
113 (2018), and Tang et al. (2019) considered tactical planning decisions including transportation, pro-
114 duction, and inventory decisions of SCND. Due to the complexity of supply chains in real-world

115 applications, Badri et al. (2013), Zokaei et al. (2017), Mönch et al. (2018), and Sagaert et al. (2019)
116 integrated strategic and tactical planning decisions in SCND problems.

117 In recent years, these SCND problems with integrated tactical and strategic planning decisions
118 have been extended in several directions. For example, reverse logistics and closed-loop supply
119 chains (Özceylan & Paksoy, 2013; Keyvanshokoo et al., 2016), finance management (Nickel et al.,
120 2012; Temocin et al., 2018; Kılıç & Uğur, 2023), pricing strategy (Azevedo et al., 2018), smart
121 warehouse operations (Zhen & Li, 2022), and supply chain resilience (Chen et al., 2023). Fattahi
122 et al. (2017) considered a multi-period SCND with the delivery lead-time of products. Mohammed
123 & Wang (2017) and Babazadeh et al. (2017) designed a three-echelon supply chain considering the
124 environmental impacts of both activities and facilities. Margolis et al. (2018) developed a multi-
125 objective model with production capacities of supply nodes for each commodity and other research,
126 considering the impact of demand on SCND problems. Gupta et al. (2021) proposed a fuzzy multi-
127 objective linear program for integrated food procurement, storage and distribution problems under
128 cost, flexibility, and quality considerations. Song et al. (2022) studied the influences of uncertainty
129 brought about by technological innovation on the supply chain in terms of supply and demand,
130 shortage penalty, and budget, and proposed a three-level model of the anti-epidemic supply chain.

131 The research about SCND, which considers the impact of demand, can be divided into three
132 categories: production is equal to demand, production is allowed to exceed demand, and production
133 is not allowed to exceed demand. The total customer demands were satisfied in research of Pan
134 & Nagi (2013), Baud-Lavigne et al. (2016), and Cheng & Tang (2018). A number of papers also
135 provide surveys of SCND problems that allow the production to exceed demand (see Sahling &
136 Kayser, 2016). Other studies do not allow production to exceed demand. For example, Guo & Li
137 (2014) ensured that the total expected ordering quantity from all selected suppliers does not exceed
138 the total expected demand and guaranteed that there was no extra inventory cost.

139 In addition, several studies pointed out that pricing greatly impacts demand. Cachon & Harker
140 (2002) indicated that price cuts would increase demand, thus reducing the average cost per unit of
141 demand. Hsieh et al. (2010) reported that pricing should be an effective tool to maximize revenue
142 when considering price-sensitive demand. Later, scholars also considered the relationship between
143 price and demand in studies of supply chain network problems. For example, Ahmadi-Javid &
144 Hoseinpour (2015b) studied an SCND problem with price-sensitive demands and inventory-capacity
145 constraints, and presented a location-inventory-pricing model. Wang et al. (2016) investigated the
146 bullwhip effect on inventory under different information sharing settings based on price-sensitive
147 demand. Ke & Bookbinder (2018) developed a tri-level programming approach to coordinate deci-
148 sions of three supply chain members on discount policies, for the case when demand is sensitive to
149 price changes. However, most scholars used demand and price as parameters when building models,
150 rather than as decision variables for direct decision making.

151 Another related line of research is the study of the SCND literature in the presence of quan-
152 tity discount. In most of the literature, studies on SCND under quantity discounts usually assume
153 that the supply chain network is single-echelon and that the problem includes only suppliers and
154 customers (Burke et al., 2008). Quantity discounting problems are usually classified into incremen-
155 tal and all-units quantity discount. Tamjizad & Mirmohammadi (2018) discussed a multi-item
156 inventory system with a discrete stochastic demand stream under continuous review, where there
157 is a limited resource and all-units quantity discount. Yang et al. (2019) proposed a deterministic
158 model to minimize the cost of service, provided that each lane being served by carriers and quantity
159 discount are taken into account. Jackson & Munson (2019) developed effective solutions for a con-
160 strained, multi-product lot-sizing problem with a common replenishment cycle, a flexible common
161 resource capacity, and all-units quantity discounts.

162 While Hsieh et al. (2010) and Ke & Bookbinder (2018) focused on price-sensitive demand and
163 quantity discounts, their research primarily delved into tactical-level decisions and did not encom-
164 pass multi-level strategic decisions. In studies similar to our present work, Ahmadi & Ghasemi
165 (2022) employed a fuzzy inference system to forecast demand and considered the impact of en-
166 ergy efficiency and environmental factors on hotel pricing to determine the optimal pricing under
167 competitive conditions. Momenitabar et al. (2022) optimized strategic and tactical decisions in a
168 closed-loop supply chain network design problem considering price-sensitive demand. However, they
169 did not directly investigate price and demand as variables and did not consider quantity discounts.
170 Safaei et al. (2022) studied a multi-level, multi-period closed-loop SCND problem by using time
171 series for demand forecasting. Their focus, however, was more on optimizing tactical-level decisions
172 and did not address price-sensitive demand and quantity discounts. Table 1 offers a comprehensive
173 overview of the key research efforts, shedding light on the primary contributions of this study.

174 Based on the literature review presented in Table 1, it is evident that in the majority of studies
175 focusing on joint optimization at the strategic and tactical levels, the optimization at the strategic
176 level primarily revolves around single-level location decision problems. The introduction of multi-
177 level networks significantly increases complexity. Furthermore, at the tactical decision level, there
178 is limited literature that simultaneously considers both price and demand as variables, with only
179 two articles addressing this aspect. Additionally, there is only one paper that takes into account
180 price-sensitive demand and quantity discounts, especially at the all-units quantity discount. Finally,
181 in the context of SCND problems, most studies tend to lean towards heuristic algorithms due to the
182 nonlinearity of the models and the complexity of solving them. Our proposed three-level network
183 structure model encompasses considerations for price-sensitive demand from the consumer side, as
184 well as quantity discounts from the supplier side. We extend this to multi-product logistics networks
185 and make decisions regarding raw material procurement through incremental quantity discount.

Table 1
Tabular literature.

Articles	OBJ	Strategic decisions				Tactical decisions					Factors			Methods
		Manu-se	Ware-se	Quan-dis	Prod	Inven	Trans	Price	Dem	Others	Multi-prod	P-D	Quan-dis	
Tamjidzad & Mir-mohammadi (2018)	CM			✓		✓					✓	All		Local search
Fattahi et al. (2015)	PM	✓			✓	✓	✓				✓			SA
Yang et al. (2019)	CM			✓			✓					All		IA
Jackson & Munson (2019)	CM							✓			✓	All		CRCA-RP
Safaei et al. (2022)	CM				✓	✓	✓		✓		✓			GA
Momenitabar et al. (2022)	CM	✓			✓	✓	✓		✓		✓			MCGP-UF
Wang et al. (2016)	CM				✓	✓			✓		✓			MMSE
Ahmadi-Javid & Hoseinpour (2015a)	PM		✓			✓			✓		✓			LRA
Ahmadi-Javid & Ghandali (2014)	PM		✓			✓			✓		✓			LRA
Ahmadi-Javid & Hoseinpour (2015b)	PM		✓			✓			✓		✓			LRA
Xiao & Qi (2016)	PM					✓			✓		✓	All		CoM
Xiao et al. (2016)	PM					✓			✓		✓			SG, NG
Our work	PM	✓		✓	✓	✓	✓	✓	✓	✓	✓	INCR		SOCP, OA

Note 1: OBJ: Objective function, Manu-se: Manufacturer selection, Ware-se: Warehouse selection, Quan-dis: Quantity discount selection, Prod: Production, Inven: Inventory, Trans: Transportation, Dem: Demand, Multi-prod: Multiple products, P-D: Price-demand, Quan-dis: Quantity discount.

Note 2: All: All-units quantity discount, INCR: incremental quantity discounts, CM: Cost minimization, PM: Profit maximization.

Note 3: SA: Simulated annealing, IA: Iterative algorithm, CRCA-RP: Common replenishment cycle approach with refinement policy, GA: Genetic algorithm, MCGP-UF: Multi-Choice Goal Programming Approach with Utility Function, MMSE: MMSE forecasting technique, LRA: Lagrangian relaxation algorithm, CoM: Coordination mechanism, SG: Stackelberg game, NG: Nash game, SOCP: Second-order cone programming, OA: outer-approximation.

186 On this basis, we tackle the nonlinearity of the model by the SOCP method to obtain an accurate
187 solution, and address large-scale instances by the outer-approximation method. In summary, the
188 main novelties and contributions of this paper can be summarized as follows:

- 189 (1) Multi-level integrated decision-making: This paper establishes an MINLP model with the objec-
190 tive of maximizing the total supply chain profit to integrate location decisions at the strategic
191 level and production, transportation, inventory, product prices, consumer demand, and discount
192 selection decisions at the tactical level.
- 193 (2) Price-sensitive demand: This paper introduces the constraint of the price-demand function in
194 the SCND problem, describes the correlation between demand and price in customer segments,
195 and captures the elastic range of demand at different price points.
- 196 (3) Incremental quantity discount: Implemented in a stepwise manner, the discount gradually in-
197 creases as the number of purchases increases, so that customers can enjoy greater price conces-
198 sions when they purchase a larger number of products. As the number of purchases increases,
199 the discount gradually increases.
- 200 (4) Solution algorithm: Due to the nonlinear nature of the model, this study proves the existence of
201 the optimal solution of the problem and the feasibility of implementing the following two algo-
202 rithms by analyzing the nature of the problem: the SOCP method and the outer-approximation
203 method. The SOCP method can obtain optimal solutions for small- and medium-scale ex-
204 periments, while the outer-approximation method can find approximate optimal solutions in a
205 reasonable time for large-scale experiments.

206 **3. SCND model considering pricing-demand relationships**

207 *3.1. Description of the problem*

208 This paper studies a multi-period, three-echelon supply chain network structure, which consists of
209 multiple manufacturers, multiple warehouses, and a single customer group. A single customer group
210 is defined as a customer type with the same consumption concept and demand, and the demands of
211 each member in the customer group are independent. The production, transportation, inventory,
212 and sales of products over multiple periods are considered in this supply chain network, as shown
213 in Figure 1.

214 Upstream in the supply chain network, manufacturers need to purchase raw materials when
215 manufacturing products. To ensure sufficient stock of raw materials, this study assumes that these
216 manufacturers choose the order-up-to-level, that is, the manufacturer n assesses the inventory at
217 each period t to determine whether it is lower than the maximum inventory level H_n^{max} , and then
218 determines the current purchase amount of raw materials q_{nt} . If the inventory is h_{nt} at the time of
219 inspection, the order quantity is $q_{nt} = H_n^{max} - h_{nt}$. The unit purchase cost of raw materials per

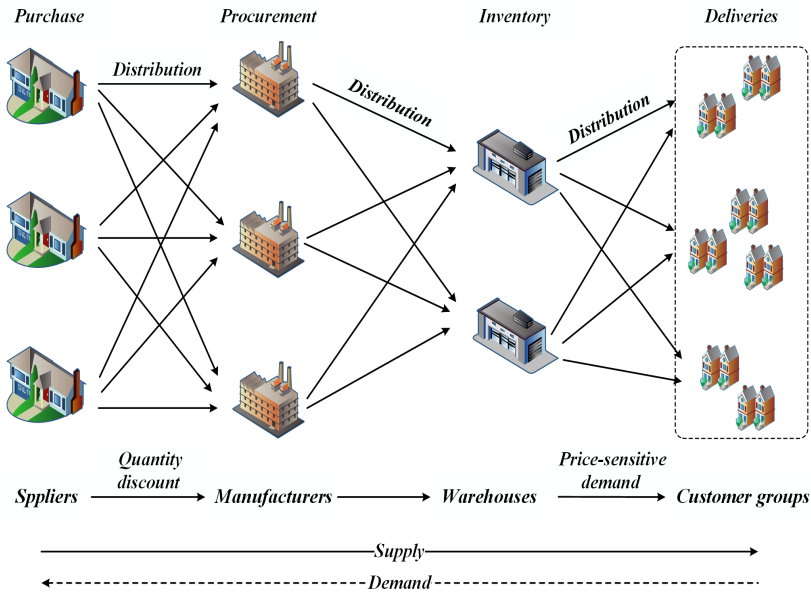


Fig. 1 Supply chain network structure

220 period is different. Meanwhile, when purchasing raw materials, manufacturers often face quantity
 221 discount provided by suppliers to reduce the procurement cost of raw materials. Figure 2 shows the
 222 cost of incremental quantity discounts corresponding to purchase quantities at different discount
 223 levels. The more the quantity purchased, the lower the price. The growth trend of the total cost
 will be slower and slower.

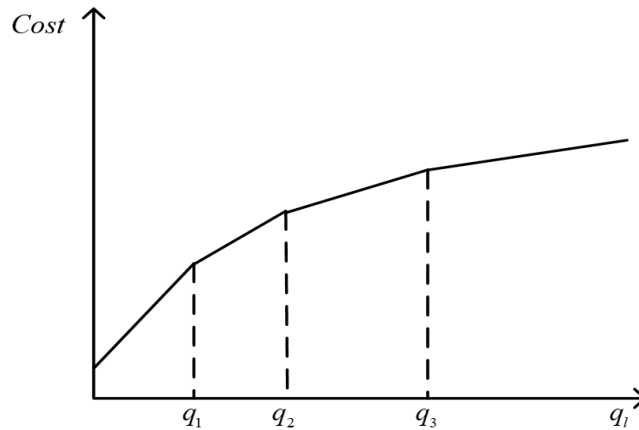


Fig. 2 Incremental quantity discount cost for order quantity q_i

224
 225 Figure 3 shows the purchase price of the incremental quantity discount corresponding to the
 226 upper limit of different discount levels. With the increase of the purchase volume, the purchase
 227 price presents a step-by-step decrease, thereby stimulating the increase of the purchase volume.

228 After the manufacturer purchases the raw materials, these purchased raw materials are used to
 229 produce the product. These produced products should be transferred to the warehouse. The sup-
 230 ply chain network contains several warehouses. Furthermore, products are transported to different

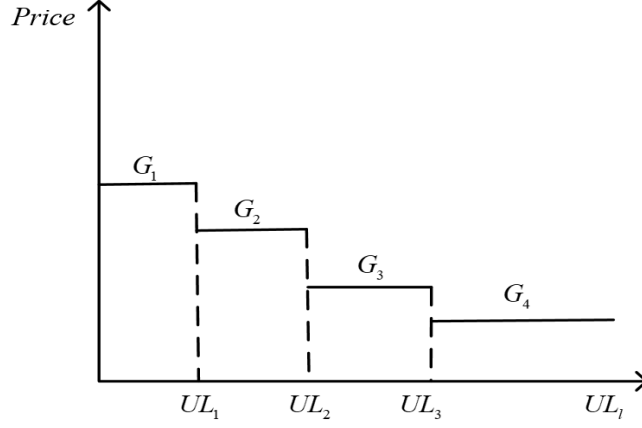


Fig. 3 Cost of purchasing raw materials for Upper limit of discount level UL_l

231 warehouses by the manufacturers' fleet, which generates transportation costs. Products shipped to
 232 the warehouses can be stored in warehouses or shipped to a customer base for sale, but products
 233 cannot be shipped between warehouses.

234 Downstream of the supply chain network, the consumer zone considers price-sensitive demand
 235 and describes the relationship between them by introducing price-demand functions. Figure 4 shows
 236 an example of the price-demand function used in this paper. When the price of product i is 0, the
 237 demand reaches the maximum D_i ; when the price of product i exceeds the consumer's expectation
 b_i , the demand is 0, i.e., no one would buy it.

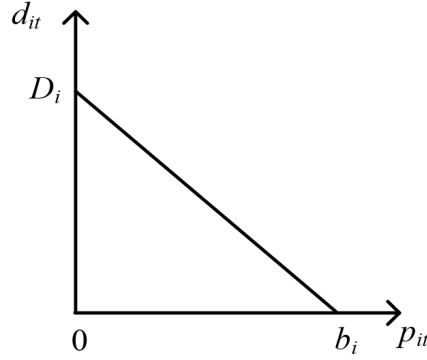


Fig. 4 Price-demand function

238

239 According to Figure 4, the functions of price and demand can be written as follows:

240

$$d_{it} = D_i - p_{it} \cdot D_i / b_i, 0 \leq p_{it} \leq b_i, \forall i \in I, t \in T. \quad (1)$$

241 3.2. Mathematical model

242 The underlying assumptions of the model are as follows:

243 (1) Manufacturers use the (t, H_n^{max}) strategy to replenish raw materials.

244 (2) The products produced by manufacturers can only be stored in warehouses.

245 The notations of this paper are as follows:

246 **Sets**

- 247 T Set of periods, $t \in T$, $T = \{1, 2, \dots, t, \dots, T\}$;
248 I Set of products, $i \in I$, $I = \{1, 2, \dots, i, \dots, I\}$;
249 W Set of warehouses, $w \in W$, $W = \{1, 2, \dots, w, \dots, W\}$;
250 N Set of manufacturers, $n \in N$, $N = \{1, 2, \dots, n, \dots, N\}$;
251 L Set of discount levels, $l \in L$, $L = \{0, 1, 2, \dots, l, \dots, L\}$;

252 **Parameters**

- 253 UL_{nl} Upper limit for discount level l for manufacturer n ;
254 F_n Fixed cost for the selection of manufacturer n ;
255 R_w Fixed cost for the opening of warehouse w ;
256 G_{nlt} Cost of purchasing one unit of raw materials by manufacturer n at discount level l during
257 period t ;
258 U_{nlt} Cost of producing one unit of product i by manufacturer n during period t ;
259 L_{nt} Cost of holding one unit of raw materials by manufacturer n during period t ;
260 L_{wit} Cost of holding one unit of product i at warehouse w during period t ;
261 V_{nwit} Cost of transporting one unit of product i between manufacturer n and warehouse
262 w during period t ;
263 C_{wit} Cost of transporting one unit of product i between warehouse w and the customer
264 zone during period t ;
265 A_i Amount of raw materials required per unit of product i ;
266 h_{n0} Amount of raw materials held at manufacturer n at the beginning of the period;
267 h_{wi0} Amount of product i held at warehouse w at the beginning of the period;
268 H_n^{max} Maximum stock level of raw materials at the manufacturer n ;
269 Ω_n^{max} Maximum production capacity at manufacturer n ;
270 K_n^{max} Maximum transportation capacity amount of manufacturer n for forwarding products to
271 all warehouses per period;
272 H_w^{max} Maximum inventory capacity at warehouse w ;
273 K_w^{max} Maximum transportation capacity amount of warehouse w for forwarding products to all
274 customers per period;
275 b_i The parameter of the price-demand function; when the price of product i exceeds b_i , the
276 demand is 0;
277 D_i Upper bound value for the demand of product i ;

278 **Decision variables**

- 279 p_{it} Price of product i during period t ;
280 d_{it} Demand of product i after its price is determined during period t ;

281	σ_{it}	Actual amount of product i purchased during period t ;
282	P_{it}	Total revenue of product i during period t ;
283	x_n	1 if manufacturer n is selected, and 0 otherwise;
284	y_w	1 if warehouse w is open, and 0 otherwise;
285	z_{nw}	1 if manufacturer n and warehouse w are selected and open, and 0 otherwise;
286	ω_{nlt}	1 if raw materials purchased by manufacturer n at discount level l during period t , and 0 otherwise;
287		
288	q_{nlt}	Amount of raw materials purchased by manufacturer n at discount level l during period t ;
289	\hat{q}_{nit}	Amount of product i produced at manufacturer n during period t ;
290	h_{nt}	Amount of raw materials held at manufacturer n during period t ;
291	h_{wit}	Amount of product i held at warehouse w during period t ;
292	λ_{nwit}	Amount of product i transported from manufacturer n to warehouse w during period t ;
293	s_{wit}	Amount of product i transported from warehouse w to customers during period t .

294 The [MINLP] mathematical model is as follows:

295 [MINLP]

$$\begin{aligned}
296 \quad \text{maximize} \quad Z = & \sum_{i \in I} \sum_{t \in T} p_{it} \sigma_{it} - \sum_{n \in N} F_n x_n - \sum_{w \in W} R_w y_w - \sum_{n \in N} \sum_{t \in T} \sum_{l \in L} G_{nlt} q_{nlt} \\
& - \sum_{n \in N} \sum_{i \in I} \sum_{t \in T} U_{nit} \hat{q}_{nit} - \sum_{n \in N} \sum_{t \in T} L_{nt} h_{nt} - \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} L_{wit} h_{wit} \\
& - \sum_{n \in N} \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} V_{nwit} \lambda_{nwit} - \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} C_{wit} s_{wit}
\end{aligned} \tag{2}$$

297 subject to

$$\sum_{l \in L} q_{nlt} = H_n^{max} - h_{nt}, \forall n \in N, t \in T, \tag{3}$$

$$q_{nlt} \geq (UL_{nl} - UL_{n(l-1)}) \omega_{n(l+1)t}, \forall n \in N, l \in \{1, \dots, L-1\}, t \in T, \tag{4}$$

$$q_{nlt} \leq (UL_{nl} - UL_{n(l-1)}) \omega_{nlt}, \forall n \in N, l \in \{1, \dots, L\}, t \in T, \tag{5}$$

$$\omega_{n(l+1)t} \leq \omega_{nlt}, \forall n \in N, l \in \{0, \dots, L-1\}, t \in T, \tag{6}$$

$$\sum_{w \in W} \lambda_{nwit} = \hat{q}_{nit}, \forall n \in N, i \in I, t \in T, \tag{7}$$

$$\sum_{i \in I} \hat{q}_{nit} \leq \Omega_n^{max} x_n, \forall n \in N, t \in T, \tag{8}$$

$$\sum_{i \in I} \lambda_{nwit} \leq K_n^{max} \cdot x_n \cdot y_w, \forall n \in N, w \in W, t \in T, \tag{9}$$

$$\sum_{i \in I} h_{wit} \leq H_w^{max} \cdot y_w, \forall w \in W, t \in T, \quad (10)$$

$$\sum_{i \in I} s_{wit} \leq K_w^{max} \cdot y_w, \forall w \in W, t \in T, \quad (11)$$

$$\sum_{w \in W} s_{wit} = \sigma_{it}, \forall i \in I, t \in T, \quad (12)$$

$$\sigma_{it} \leq d_{it}, \forall i \in I, t \in T, \quad (13)$$

$$d_{it} = D_i - p_{it} \cdot D_i / b_i, 0 \leq p_{it} \leq b_i, \forall i \in I, t \in T, \quad (14)$$

$$h_{n(t-1)} + \sum_{l \in L} q_{nlt} = h_{nt} + \sum_{i \in I} A_i \cdot \hat{q}_{nit}, \forall n \in N, t \in T \setminus \{1\}, \quad (15)$$

$$h_{n0} + \sum_{l \in L} q_{nl1} = h_{n1} + \sum_{i \in I} A_i \cdot \hat{q}_{ni1}, \forall n \in N, \quad (16)$$

$$h_{wi(t-1)} + \sum_{n \in N} \lambda_{nwit} = h_{wit} + s_{wit}, \forall w \in W, i \in I, t \in T \setminus \{1\}, \quad (17)$$

$$h_{wi0} + \sum_{n \in N} \lambda_{nwit} = h_{wit} + s_{wit}, \forall w \in W, i \in I, t = 1, \quad (18)$$

$$x_n, y_w, z_{nw}, \omega_{nlt} \in \{0, 1\}, \forall n \in N, w \in W, t \in T, \quad (19)$$

$$p_{it}, d_{it}, q_{nlt}, \hat{q}_{nit}, h_{nt}, h_{wit}, \lambda_{nwit}, s_{wit} \geq 0, \forall n \in N, w \in W, l \in L, i \in I, t \in T, \quad (20)$$

$$\sigma_{it} \geq 0, \forall i \in I, t \in T. \quad (21)$$

298 The objective function (2) describes the total profit of the supply chain network. The first
 299 item in the objective function represents the total revenue of the products sold; the second item
 300 represents the construction cost of warehouses; the third item represents the construction cost of
 301 manufacturers; the fourth item represents the procurement cost of raw materials; the fifth item
 302 represents the production cost at manufacturers; the sixth item represents the inventory cost of the
 303 raw materials at manufacturers; the seventh item represents the inventory cost of products at the
 304 warehouses; and the last two items represent the transportation cost of the products. Constraints
 305 (3) represent the purchase amount of raw materials at the beginning of period t . Constraints (4) and
 306 Constraints (5) ensure that the purchase quantity of raw materials at each discount level is between

307 the lower and upper limits according to the incremental quantity discount of each supplier in each
308 period. Constraints (6) help determine the discount level for raw materials purchased from each
309 supplier in each period. Constraints (7) indicate that all products produced by manufacturer n are
310 shipped to warehouse w . Constraints (8) are capacity constraints, indicating that the manufacturers'
311 production levels cannot exceed their maximum production capacity. Constraints (9) indicate that
312 the transportation volume of the products from the manufacturers to the warehouses cannot exceed
313 the maximum transportation capacity of manufacturers. Constraints (10) indicate that the inventory
314 of the products in warehouse w cannot exceed its maximum inventory capacity. Constraints (11)
315 indicate that the transportation volume of the products from warehouse w to the consumer cannot
316 exceed its maximum transportation capacity. Constraints (12) indicate that the quantity of products
317 transported from the warehouse must meet the demands of consumers. Constraints (13) indicate
318 that the actual purchase amount of consumers is less than or equal to their demand after the price
319 is determined. Constraints (14) represent the price-demand relationships. Constraints (15) and
320 (16) represent the flow balance of the products production. Constraints (17) and (18) represent the
321 flow balance of the products transportation. Constraints (19) clarify the numerical definition of 0-1
322 variables. Constraints (20) and (21) clarify the numerical definition of the non-negative variables.

323 In the model [MINLP], it is obvious that two nonlinearities exist. One is in the first item of the
324 objective function and the other is in constraints (9). Here, the decision variable z_{nw} is introduced
325 to convert constraints (9) to the following constraints:

$$\sum_{i \in I} \lambda_{nwit} \leq K_n^{max} \cdot z_{nw}, \forall n \in N, w \in W, t \in T, \quad (22)$$

$$z_{nw} \leq x_n, \forall n \in N, w \in W, \quad (23)$$

$$z_{nw} \leq y_w, \forall n \in N, w \in W, \quad (24)$$

$$z_{nw} \in \{0, 1\}, \forall n \in N, w \in W. \quad (25)$$

326 4. Solution for the SCND model considering pricing-demand relationships

327 Due to the nonlinearity of the first item in the objective function (2), solving the [MINLP] mathe-
328 matical model is particularly difficult. Discretizing the price is the most common way to deal with
329 the nonlinear part of the model. Fattahi et al. (2015) discretized the pricing level by introducing
330 the price-demand functions and solved it by using a linear relaxation-based simulated annealing
331 algorithm. However, the proposed algorithm introduces additional integer decision variables, which
332 affect the quality of the solution. Therefore, the properties of the MINLP model are first analyzed,
333 and then, two different algorithms are proposed accordingly in this research.

334 *4.1. Analysis of the SCND model considering pricing-demand relationships*

335 The model [MINLP] has the following characteristics and can derive the following relevant theorems:

336 **Theorem 1:** The model [MINLP] has an optimal solution.

337 **Proof.** All constraints in [MINLP] are linear and are connected by \leq , $=$, and \geq . There are
 338 no $<$ and $>$ constraints. Therefore, the feasible region of [MINLP] is closed. The feasible region of
 339 [MINLP] is bounded because of the constraints (3), (8)-(11), (13), (14), (20), and (21). Hence, the
 340 feasible region is a compact set. The feasible region is not empty, because there is a feasible solution
 341 where all decision variables are zero. Therefore, the model [MINLP] has an optimal solution.

342 **Theorem 2:** The decision variable $(\sigma_{it}, d_{it}, \forall i \in I, t \in T)$ of [MINLP] has an optimal solution
 343 $(\sigma_{it}^*, d_{it}^*, \forall i \in I, t \in T)$. For $\forall i \in I, t \in T$, if $\sigma_{it}^* > 0$, then $\sigma_{it}^* = d_{it}^*$.

344 **Proof.** The theorem is proved by contradiction. Theorem 1 illustrates the existence of an
 345 optimal solution for [MINLP]. Suppose that an optimal solution $(\hat{\sigma}_{it}, \hat{d}_{it}, \forall i \in I, t \in T)$ exists and
 346 the solution satisfies $0 < \hat{\sigma}_{i\bar{t}} < \hat{d}_{i\bar{t}}$ for $\bar{i} \in I, \bar{t} \in T$.

347 Since $p_{i\bar{t}}(d_{i\bar{t}})$ is strictly decreasing on $[0, D_i]$ and $0 < \hat{\sigma}_{i\bar{t}} < \hat{d}_{i\bar{t}} \leq D_i$, then $p_{i\bar{t}}(\sigma_{i\bar{t}}) > p_{i\bar{t}}(d_{i\bar{t}})$.
 348 We have $p_{i\bar{t}}(\sigma_{i\bar{t}})\sigma_{i\bar{t}} > p_{i\bar{t}}(d_{i\bar{t}})\sigma_{i\bar{t}}$ because $\sigma_{i\bar{t}} > 0$. Therefore, if the price $p_{i\bar{t}}(d_{i\bar{t}})$ increases
 349 to $p_{i\bar{t}}(\sigma_{i\bar{t}})$, the profit (2) will also increase. Then, the new solution $(\sigma_{it} = \hat{\sigma}_{it}, \forall i \in I, t \in$
 350 $T; d_{it} = \hat{d}_{it}, \forall i \in I \setminus \{\bar{i}\}, t \in T \setminus \{\bar{t}\}; d_{i\bar{t}} = \hat{\sigma}_{i\bar{t}})$ is feasible and better than the original solution
 351 $(\hat{\sigma}_{it}, \hat{d}_{it}, \forall i \in I, t \in T)$, which contradicts the assumption that the solution $(\hat{\sigma}_{it}, \hat{d}_{it}, \forall i \in I, t \in T)$
 352 is optimal. Theorem 2 is proved.

353 Therefore, constraints (12), (13), and (14) can be transferred to:

354
$$\sum_{w \in W} s_{wit} = D_i - p_{it} \cdot D_i/b_i, 0 \leq p_{it} \leq b, \forall i \in I, t \in T. \quad (26)$$

355 By proving the properties of the optimal solution, actual purchases σ_{it} can then be replaced by
 356 demand d_{it} in the objective function (2).

357 *4.2. Second-order cone programming*

358 The nonlinearity in the first term of the objective function cannot be solved directly using the
 359 mathematical programming software CPLEX. Therefore, the problem was converted into an SOCP
 360 problem to obtain the exact optimal solution in this section.

361 Theorem 3 can be obtained from the introduction of the price-demand functions, see constraints
 362 (1) and Figure 3.

363 **Theorem 3:** The function $d_{it}(p_{it})$ is strictly decreasing on $[0, b_i]$ and conforms to the Lipschitz
 364 continuity.

365 The function $d_{it}(p_{it})$ has an inverse function due to strict monotonicity, and the inverse function
 366 can be expressed as:

367

$$p_{it} = b_i - d_{it} \cdot b_i/D_i, 0 \leq d_{it} \leq D_i, \forall i \in I, t \in T. \quad (27)$$

368 According to constraints (27), the function $p_{it}(d_{it})$ is also a strictly decreasing function as shown
in Figure 4.

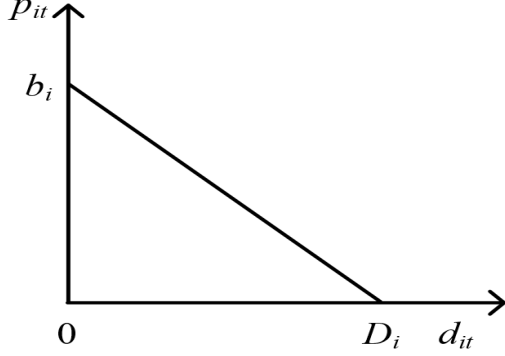


Fig. 5 Demand-price function

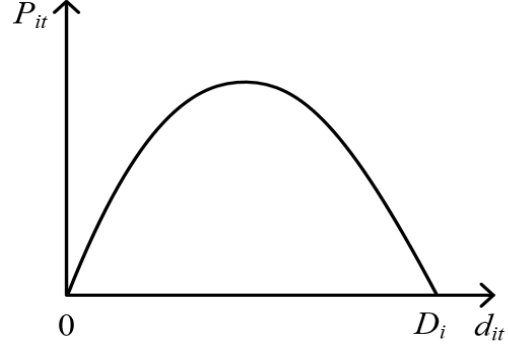


Fig. 6 Demand-revenue function

369

370 The rationality of the Lipschitz continuity is that demands originate from multiple consumers.
371 When the price changes a little, the demand does not fluctuate significantly.

372 P_{it} is defined as the revenue of product i during period t , i.e., $P_{it} = p_{it} \cdot d_{it}$. According to
373 constraints (27), we can get:

374

$$P_{it} = (-d_{it}^2 + D_i \cdot d_{it}) \cdot b_i/D_i, 0 \leq d_{it} \leq D_i, \forall i \in I, t \in T. \quad (28)$$

375 **Theorem 4:** P_{it} is strictly concave on $[0, D_i]$ and continuously differentiable about d_{it} .

376 The demand-revenue function is shown in Figure 5.

377 Therefore, constraints (28) can be transferred into constraints (29):

378

$$P_{it} \leq (-d_{it}^2 + D_i \cdot d_{it}) \cdot b_i/D_i, 0 \leq d_{it} \leq D_i, \forall i \in I, t \in T. \quad (29)$$

379 Furthermore, constraints (29) can be re-written as:

380

$$-d_{it}^2 + D_i \cdot d_{it} \geq D_i \cdot P_{it}/b_i, 0 \leq d_{it} \leq D_i, \forall i \in I, t \in T, \quad (30)$$

$$d_{it}^2 - D_i \cdot d_{it} \leq -D_i \cdot P_{it}/b_i, 0 \leq d_{it} \leq D_i, \forall i \in I, t \in T. \quad (31)$$

381 Then, constraints (31) can be transformed to a set of hyperbolic inequalities, as shown in the
382 following:

383

$$(d_{it} - D_i/2)^2 \leq -D_i \cdot P_{it}/b_i + D_i^2/4, 0 \leq d_{it} \leq D_i, \forall i \in I, t \in T. \quad (32)$$

384 The above hyperbolic inequalities can be rewritten as a set of SOCP constraints:

$$385 \quad \|(2d_{it} - D_i, -D_i \cdot P_{it}/b_i + D_i^2/4 - 1)\|_2 \leq -D_i \cdot P_{it}/b_i + D_i^2/4 + 1, 0 \leq d_{it} \leq D_i, \forall i \in I, t \in T, \quad (33)$$

386 where $\|\cdot\|_2$ denotes the Euclidean norm.

387 Therefore, the model [MINLP] can be changed to model [MINLP-SOCP]:

388 [MINLP-SOCP]

$$389 \quad \begin{aligned} \text{maximize } Z = & \sum_{i \in I} \sum_{t \in T} P_{it} - \sum_{n \in N} F_n x_n - \sum_{w \in W} R_w y_w - \sum_{n \in N} \sum_{l \in L} \sum_{t \in T} G_{nlt} q_{nlt} \\ & - \sum_{n \in N} \sum_{i \in I} \sum_{t \in T} U_{nit} \hat{q}_{nit} - \sum_{n \in N} \sum_{t \in T} L_{nt} h_{nt} - \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} L_{wit} h_{wit} \\ & - \sum_{n \in N} \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} V_{nwit} \lambda_{nwit} - \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} C_{wit} s_{wit} \end{aligned} \quad (34)$$

390 subject to constraints (3)-(8), (10), (11), (15)-(20), (22)-(26), and (33).

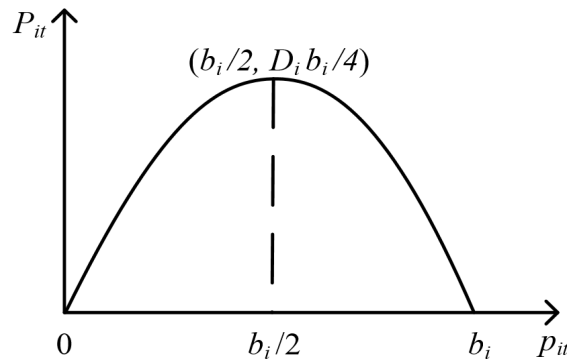
391 4.3. Outer-approximation

392 Although the SOCP model can obtain the exact optimal solution, it can only deal with small and
393 medium-sized problems in a reasonable time. Therefore, the outer-approximation method is pro-
394 posed in this section to find the approximate optimal solution of the problem and improve the
395 efficiency of solving large-scale problems.

396 As mentioned above, define P_{it} as the revenue of product i during period t , and the revenue
397 function changes from $P_{it} = p_{it} \cdot \sigma_{it}$ to $P_{it} = p_{it} \cdot d_{it}$ by the above theorems. According to constraints
398 (14), the following constraints can be obtained:

$$399 \quad P_{it} = D_i \cdot p_{it} - p_{it}^2 \cdot D_i/b_i, 0 \leq p_{it} \leq b_i, \forall i \in I, t \in T. \quad (35)$$

400 According to constraints (35), the maximum value of P_t on $[0, b_i]$ is obtained at $b_i/2$ (symmetry
401 axis) and the maximum value is $D_i \cdot b_i/4$. The graph of the function $P_{it}(p_{it})$ on $[0, b_i]$ is shown in
Figure 6.



402 **Fig. 7** Price-revenue function

403 The model [MINLP] can be transferred to the [MINLP-OA1] model (where OA stands for outer-
 404 approximation).

405 [MINLP-OA1]

$$\begin{aligned}
 406 \quad \text{maximize } Z = & \sum_{i \in I} \sum_{t \in T} P_{it} - \sum_{n \in N} F_n x_n - \sum_{w \in W} R_w y_w - \sum_{n \in N} \sum_{l \in L} \sum_{t \in T} G_{nlt} q_{nlt} \\
 & - \sum_{n \in N} \sum_{i \in I} \sum_{t \in T} U_{nit} \hat{q}_{nit} - \sum_{n \in N} \sum_{t \in T} L_{nt} h_{nt} - \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} L_{wit} h_{wit} \\
 & - \sum_{n \in N} \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} V_{nwit} \lambda_{nwit} - \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} C_{wit} s_{wit}
 \end{aligned} \tag{36}$$

407 subject to constraints (3)-(8), (10), (11), (15)-(20), (22)-(27), and (35).

408 Since the function $P_{it}(p_{it})$ is concave on $[0, b_i]$, the outer-approximation method is used to
 409 approximate it.

410 For $\forall p_{it} \in [0, b_i]$, P_{it} is approximated by using the sum of many piecewise-linear functions
 411 and its approximation error e can be controlled within a predetermined tolerance level. Define
 412 $\Psi = \{1, 2, \dots, |\Psi|\}$ as the set of tangents and the tangent's expression is: $r_{it}^\psi \times p_{it} + sk_{it}^\psi$, where r_{it}^ψ
 413 and sk_{it}^ψ represent the slope and intercept of tangent ψ , respectively. Then, the piecewise-linear
 414 approximation functions can be written as:

$$415 \quad \bar{P}_{it} = \min \left\{ r_{it}^\psi \times p_{it} + sk_{it}^\psi, \forall i \in I, t \in T, \psi \in \Psi \right\}. \tag{37}$$

416 The characteristics of the piecewise-linear function $\bar{P}_{it}(p_{it})$ and the concave function $P_{it}(p_{it})$
 417 indicate:

$$418 \quad \bar{P}_{it} \geq P_{it}, \forall i \in I, t \in T. \tag{38}$$

419 The following constraints can be obtained by constraints (37) and (38):

$$420 \quad P_{it} \leq r_{it}^\psi \times p_{it} + sk_{it}^\psi, \forall i \in I, t \in T, \psi \in \Psi. \tag{39}$$

421 Therefore, the model [MINLP-OA1] can be modified as model [MINLP-OA2]:

422 [MINLP-OA2]

$$\begin{aligned}
 423 \quad \text{maximize } Z = & \sum_{i \in I} \sum_{t \in T} P_{it} - \sum_{n \in N} F_n x_n - \sum_{w \in W} R_w y_w - \sum_{n \in N} \sum_{t \in T} \sum_{l \in L} G_{nlt} q_{nlt} \\
 & - \sum_{n \in N} \sum_{i \in I} \sum_{t \in T} U_{nit} \hat{q}_{nit} - \sum_{n \in N} \sum_{t \in T} L_{nt} h_{nt} - \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} L_{wit} h_{wit} \\
 & - \sum_{n \in N} \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} V_{nwit} \lambda_{nwit} - \sum_{w \in W} \sum_{i \in I} \sum_{t \in T} C_{wit} s_{wit}
 \end{aligned} \tag{40}$$

424 subject to constraints (3)-(8), (10), (11), (15)-(20), (22)-(26), and (39).

425 To make it easier to understand, the tangent generating steps of the piecewise-linear approxima-
 426 tion functions are described in the following. Figure 7 illustrates the algorithm (taking five tangent
 427 lines of period t as an example).

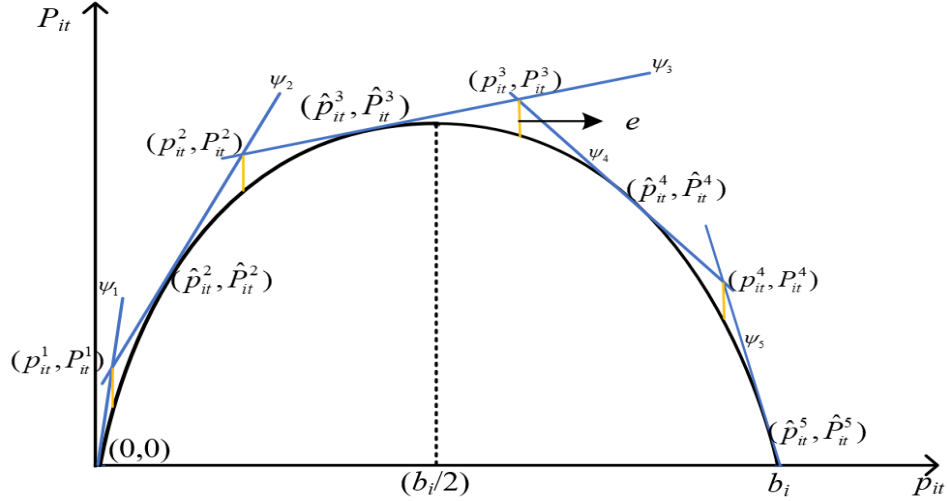


Fig. 8 Generation of the piecewise-linear approximation functions

428 Step 0: Set $\psi = 1$ and the tangent point $(\hat{p}_{it}^1, \hat{P}_{it}^1) = (0, 0)$. Pass the point $(\hat{p}_{it}^1, \hat{P}_{it}^1)$ and draw
 429 the tangent ψ of the function $P_{it}(p_{it})$.

430 Step 1: Calculate the slope r_{it}^ψ and intercept sk_{it}^ψ of the tangent ψ :

431
$$r_{it}^\psi = D_i \cdot (b_i - 2\hat{p}_{it}^\psi) / b_i, \quad (41)$$

$$sk_{it}^\psi = \hat{P}_{it}^\psi - r_{it}^\psi \cdot \hat{p}_{it}^\psi. \quad (42)$$

432 Step 2: Set the error e between the tangent ψ and the function $P_{it}(p_{it})$. At the error value
 433 point, there is $P_{it}(p_{it}) = \psi - e$. We can then obtain the error value point $(p_{it}^\psi, P_{it}^\psi)$:

434
$$P_{it}^\psi = r_{it}^\psi \cdot p_{it}^\psi + sk_{it}^\psi. \quad (43)$$

435 Step 3: Determine whether $p_{it}^\psi \leq b_i$. If yes, let $\psi = \psi + 1$ and continue at Step 4. Otherwise,
 436 the calculations stay ends and go to Step 5.

437 Step 4: Pass the point $(p_{it}^{\psi-1}, P_{it}^{\psi-1})$ and draw the tangent ψ . Go to Step 1.

438 Step 5: Obtain all the tangent lines.

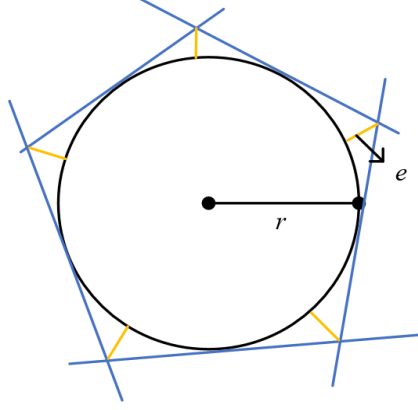
439 Based on the outer-approximation of the revenue function $P_{it}(p_{it})$, the model can be directly
 440 solved by adding linear constraints (39).

441 **Theorem 5:** The lower bound for the number of tangent lines to the revenue function $P_{it}(p_{it})$
 442 is $2\pi / \left\{ \arccos \left[\sqrt{(1 + D_i^2)^3} / \left(\sqrt{(1 + D_i^2)^3} - 2D_i e / b_i \right) \right] \right\}$, and the upper bound for the number
 443 is $2\pi / \arccos [-b_i / (2D_i e)]$.

444 **Proof.** The method for calculating the number of tangent lines to the parabola function $P_{it}(p_{it})$
 445 is:

446 At first, we take a circle with circumference l and radius r (as shown in Figure 8) as an example.
 447 The circle can be approximated by a group of tangent lines. Furthermore, the number of the tangent

448 lines is decided by the the error value e . The curvature of a circle is fixed, whereas the curvature
 449 of a parabola constantly changes. Therefore, the number of tangent lines can be calculated by the
 parabolic curvature.



450 **Fig. 9** Schematic diagram of the number of required tangent lines

451 Furthermore, it is obviously that the greater the curvature is, the more the tangents are required.
 452 So the upper bound of the number of tangents can be obtained at the maximum curvature. Similarly,
 453 the lower bound of the number of tangents can be obtained at the minimum curvature.

454 Then, we can calculate the number of tangent lines of the revenue function $P_{it}(p_{it})$ as follows:

455 The curvature radius of the parabola $y = ax^2 + bx + c(a \neq 0)$ is $\rho = (1 + y')^3 / y''$. The
 456 curvature radius at the point $(-b/2a, (4ac - b^2)/4a)$ is $1/2a$, and the curvature radius at $(-b/a, 0)$
 457 is $\sqrt{(1 + b^2)^3} / 2a$.

458 (1) When the curvature radius is $1/2a$, the maximum curvature is obtained and the circumference
 459 of a circle is π/a . Then the upper bound for the number of tangent lines can be calculated.
 460 When the error is e , the length of the circumference that a tangent can cut is $\arccos(2ae)^{-1}/2a$,
 461 and the number of required tangent lines is $2\pi / \arccos(2ae)^{-1}$.

462 (2) When the curvature radius is $\sqrt{(1 + b^2)^3} / 2a$, the minimum curvature is obtained and the cir-
 463 cumference of a circle is $\pi\sqrt{(1 + b^2)^3} / a$. Then the lower bound for the number of tangent lines
 464 can be calculated. When the error is e , the length of the circumference that can be cut by a tan-
 465 gent is $\sqrt{(1 + b^2)^3} \cdot \arccos \left[\frac{\sqrt{(1 + b^2)^3}}{\left(\sqrt{(1 + b^2)^3} + 2ae\right)} \right] / 2a$, and the number of required
 466 tangent lines is $2\pi / \left\{ \arccos \left[\frac{\sqrt{(1 + b^2)^3}}{\left(\sqrt{(1 + b^2)^3} - 2ae\right)} \right] \right\}$.

467 By putting the parameters of constraints (35) into the above conclusions, both the maximum
 468 and minimum curvature radii of the revenue function $P_{it}(p_{it})$ can be thus obtained. Then, the
 469 lower bound (ψ_{LB}) and upper bound (ψ_{UB}) for the numbers of tangent lines to the function $P_{it}(p_{it})$
 470 can be calculated: $\psi_{LB} = 2\pi / \left\{ \arccos \left[\frac{\sqrt{(1 + D_i^2)^3}}{\left(\sqrt{(1 + D_i^2)^3} - 2D_i e / b_i\right)} \right] \right\}$ and $\psi_{UB} =$
 471 $2\pi / \arccos[-b_i / (2D_i e)]$.

472 Theorem 3 is proved.

473 5. Numerical experiments

474 To assess the performance of the model and the algorithms proposed in this paper, the mathematical
 475 programming software CPLEX (12.5.1) was used to compare solutions solved by different algorithms
 476 at the same scale. These numerical experiments were conducted on a personal computer with Intel
 477 Core i7-4790 CPU (3.60 GHz), with 16.00 GB of RAM.

478 5.1. Results for three different scale problems

479 The relevant parameters involved in the experiments are uniformly distributed within the range
 480 shown in Table 2. For example, the values of F_n (Fixed cost for the selection of manufacturer n)
 481 are randomly distributed in [7000, 8000]. The experiment uses 3 discount levels. The parameters
 482 related to the discount level are uniformly distributed in the ranges given in Table 3.

Table 2
Parameters setting for the experiments.

Parameters	Value	Parameters	Value	Parameters	Value
F_n	[7000, 8000]	L_{wit}	[0.5, 1.5]	U_{nit}	[5, 6.5]
R_w	[3000, 4000]	L_{nt}	[0.3, 1.3]	C_{wit}	[1.2, 2.2]
483 D_i	[4000, 5000]	V_{nwit}	[1, 2]	h_{wi0}	[40, 60]
b_i	[50, 55]	h_{n0}	[100, 150]	Ω_n^{max}	[1500, 2200]
A_i	[3, 5]	H_n^{max}	[2000, 2200]	K_w^{max}	[1000, 1200]
K_n^{max}	[2200, 2500]	K_w^{max}	[800, 1000]		

Table 3
Parameter settings related to different discount levels.

The discount level	Parameters	Value	Parameters	Value
1	G_{n1t}	[10, 13]	UL_{n1}	[800, 1000]
2	G_{n2t}	[9, 10]	UL_{n2}	[1800, 2000]
3	G_{n3t}	[8, 9]	UL_{n3}	[2800, 3000]

484
 485 First, a sensitivity analysis experiment of the error value e was conducted to select the
 486 appropriate parameter for the outer-approximation algorithm. Table 4 shows the effects of different
 487 error values e on the solution results for $N = 3$, $W = 9$, $I = 3$, and $T = 10$. Table 4 shows that the
 488 gaps of the objective values between the outer-approximation algorithm and the SOCP algorithm
 489 decrease with gradually decreasing error value e . However, the computing time increases with
 490 decreasing error value e . The gap between both algorithms is relatively small and the computing
 491 time is relatively short at an error value of 5. Therefore, an error value of 5 ($e = 5$) was chosen for
 492 the outer-approximation algorithm to compare the solution quality and CPU time between both
 493 algorithms.

Table 4
Error value sensitivity analysis.

Error Values	Time (s)		<i>OBJ</i>		GAP _{OBJ} (%)
	OA	SOCP	OA	SOCP	
50	2.265		185897		0.794
40	2.706		185569		0.616
30	2.808		185246		0.441
20	2.859		184921		0.265
10	3.941	8.775	184818	184433	0.209
5	4.592		184654		0.120
1	10.352		184597		0.089
0.1	60.288		184456		0.012
0.01	90.067		184443		0.005

Notes: (1) Time is the computation time (s). (2) *OBJ* shows objective values. (3) GAP_{OBJ} is calculated by: $(OBJ_{OA} - OBJ_{SOCP})/OBJ_{SOCP}$.

494 To ensure the validity of the comparison between both algorithms, this study sets experimental
 495 parameters for three sets of different scales, including small-, medium-, and large-sized experiments.
 496 Tables 5, 6, and 7 show the numerical experimental results of both the outer-approximation
 algorithm and the SOCP algorithm in three scales.

Table 5
Comparison of experimental results of both developed algorithms for small scales.

Instances	<i>T</i>	<i>N</i>	<i>W</i>	<i>I</i>	Time (s)		<i>OBJ</i>		GAP _{OBJ} (%)
					OA	SOCP	OA	SOCP	
1	5	1	5	2	0.651	0.865	123174	123110	0.052
2	5	1	5	3	0.605	0.648	191598	191432	0.087
3	5	1	7	3	0.596	0.714	198030	197873	0.080
4	5	1	7	4	0.828	1.171	265531	265381	0.057
5	5	2	7	4	0.863	1.190	275589	275448	0.051
6	5	2	7	5	0.688	1.050	345659	345498	0.047
7	5	2	9	5	0.940	1.696	377950	377815	0.036
8	5	2	9	6	1.155	3.534	437017	436757	0.060

Notes: (1) Time represents the computation time (s). (2) *OBJ* represents the objective values. (3) GAP_{OBJ} is calculated by: $(OBJ_{OA} - OBJ_{SOCP})/OBJ_{SOCP}$.

497
 498 Since the SOCP algorithm can obtain an optimal solution, the gap and CPU time were
 499 compared between the outer-approximation algorithm and the SOCP algorithm. Table 5 shows
 500 that the outer-approximation algorithm and the SOCP algorithm can obtain the solution in a
 501 shorter computation time for small-scale problems, and the gap of results between both algorithms
 502 is very small. The average gap value is about 0.0588% for the outer-approximation algorithm and
 503 the SOCP algorithm.

504 The gaps of results are also not evident for medium-scale instances in Table 6, and the average
 505 gap value is about 0.1144% for the outer-approximation algorithm and the SOCP algorithm.

Table 6
Comparison of experimental results of both developed algorithms for middle scales.

Instances	T	N	W	I	Time (s)		OBJ		GAP $_{OBJ}$ (%)
					OA	SOCP	OA	SOCP	
9	10	3	9	3	55.365	53.903	234684	233512	0.502
10	10	3	9	4	50.758	120.534	336636	336352	0.084
11	10	3	11	4	35.934	30.206	350508	350206	0.086
12	10	3	11	5	16.085	19.333	487126	486835	0.060
13	10	4	11	5	26.832	217.654	489126	488848	0.057
14	10	4	11	6	45.431	1131.526	569665	569189	0.084
15	10	4	13	6	51.661	1117.336	589043	587256	0.304
16	10	4	13	7	176.904	3296.290	674985	674122	0.128

Notes: (1) Time represents the computation time (s). (2) OBJ represents the objective values. (3) GAP $_{OBJ}$ is calculated by: $(OBJ_{OA} - OBJ_{SOCP})/OBJ_{SOCP}$.

506 However, the CPU time of the SOCP algorithm increased noticeably compared to the outer-
507 approximation algorithm.

508 Table 7 also provides a comparison between the outer-approximation algorithm and the SOCP
509 algorithm for large-scale instances. For the SOCP algorithm, the optimal solution cannot be
510 obtained for more than 4 hours. Experiments indicate that the SOCP algorithm cannot solve
511 instances where the scale of the case exceeds 20 periods, 5 manufacturers, 13 warehouses, and
512 4 products; however, the outer-approximation algorithm can solve these cases within reasonable
computing time.

Table 7
Comparison of experimental results of both developed algorithms for large scales.

Instances	T	N	W	I	Time (h)		OBJ		GAP $_{OBJ}$ (%)
					OA	SOCP	OA	SOCP(*)	
17	20	5	13	4	0.9	≥ 4	315875	298542	-
18	20	5	13	5	1.1	≥ 4	594788	575853	-
19	20	5	15	5	1.5	≥ 4	680302	659475	-
20	20	5	15	6	1.8	≥ 4	751012	721217	-
21	20	6	15	6	1.8	≥ 4	708496	669917	-
22	20	6	15	7	1.9	≥ 4	852931	819432	-
23	20	6	17	7	2.1	≥ 4	899782	850263	-
24	20	6	17	8	2.6	≥ 4	1102087	1089521	-

Notes: (1) Time represents the computation time (h). (2) OBJ represents the objective values. (3) GAP $_{OBJ}$ is calculated by: $(OBJ_{OA} - OBJ_{SOCP})/OBJ_{SOCP}$. (4) ≥ 4 indicates that the computation time is more than 4 h and the optimal solution cannot be found. (5) (*) indicates that the value is the result of the experimental calculation of 4 h. (6) (-) represents that the value cannot be calculated.

513

514 Based on the numerical results, the following conclusions can be drawn:

- 515 (1) For small and medium-sized experiments, SOCP algorithm can obtain the optimal solution in a
516 short time. the outer-approximation algorithm can also find the approximate optimal solution in
517 an effective time. In terms of solving accuracy, the average gap between the outer-approximation

518 method and SOCP algorithm is 0.0588%.

519 (2) For large-sized experiments, the SOCP algorithm cannot solve the model, while the outer-
 520 approximation algorithm can still obtain approximate solutions within a reasonable time. This
 521 implies that the outer-approximation algorithm is an effective method to solve the proposed
 522 MINLP model on large-sized experiments.

523 *5.2. Comparison with the results of the problems without discount levels*

524 Tables 8, 9, and 10 show the experimental results of the two algorithms at three different scales
 525 without considering the discount level. The advantage of the outer-approximation algorithm in
 526 solving medium- and large-scale problems is reflected by the results of comparing the quality and
 efficiency of the algorithms.

Table 8
 Experimental results of both developed algorithms for small scales without discount levels.

Instances	T	N	W	I	Time (s)		OBJ		GAP_{OBJ} (%)
					OA	SOCP	OA	SOCP	
1	5	1	5	2	0.192	0.151	112475	112248	0.00202
2	5	1	5	3	0.195	0.194	181536	180897	0.00353
3	5	1	7	3	0.294	0.161	192591	192353	0.00124
4	5	1	7	4	0.280	0.170	252403	251959	0.00176
5	5	2	7	4	1.244	1.186	247303	247040	0.00106
6	5	2	7	5	0.397	0.197	324215	323884	0.00102
7	5	2	9	5	0.348	0.208	357686	357149	0.00150
8	5	2	9	6	0.685	0.589	413947	413488	0.00111

Notes: (1) Time represents the computation time (s). (2) OBJ represents the objective values. (3) GAP_{OBJ} is calculated by: $(OBJ_{OA} - OBJ_{SOCP})/OBJ_{SOCP}$.

Table 9
 Experimental results of both developed algorithms for middle scales without discount levels.

Instances	T	N	W	I	Time (s)		OBJ		GAP_{OBJ} (%)
					OA	SOCP	OA	SOCP	
9	10	3	9	3	10.226	7.175	185105	184433	0.00364
10	10	3	9	4	11.963	30.549	265182	264502	0.00257
11	10	3	11	4	3.967	0.447	263976	263300	0.00257
12	10	3	11	5	3.560	20.001	393326	392635	0.00176
13	10	4	11	5	28.458	0.598	377904	377219	0.00182
14	10	4	11	6	21.113	78.755	461111	460631	0.00104
15	10	4	13	6	8.423	211.062	474663	473702	0.00203
16	10	4	13	7	13.587	2418.346	568995	568252	0.00131

Notes: (1) Time represents the computation time (s). (2) OBJ represents the objective values. (3) GAP_{OBJ} is calculated by: $(OBJ_{OA} - OBJ_{SOCP})/OBJ_{SOCP}$.

527

528 Table 11 compares the results of the two models obtained by the SOCP and the outer-
 529 approximation method at different scales. We can see that considering the incremental quantity

Table 10

Experimental results of both developed algorithms for large scales without discount levels.

Instances	T	N	W	I	Time (s)		OBJ		GAP_{OBJ} (%)
					OA	SOCP	OA	SOCP(*)	
17	20	5	13	4	394.585	$\geq 4h$	238495	229156	-
18	20	5	13	5	644.393	$\geq 4h$	444029	435056	-
19	20	5	15	5	965.325	$\geq 4h$	503654	499863	-
20	20	5	15	6	1940.635	$\geq 4h$	557763	549765	-
21	20	6	15	6	2399.359	$\geq 4h$	519251	513265	-
22	20	6	15	7	2965.354	$\geq 4h$	649691	642651	-
23	20	6	17	7	2657.422	$\geq 4h$	664633	656210	-
24	20	6	17	8	3448.729	$\geq 4h$	798300	790280	-

Notes: (1) Time represents the computation time (s). (2) OBJ represents the objective values. (3) GAP_{OBJ} is calculated by: $(OBJ_{OA} - OBJ_{SOCP})/OBJ_{SOCP}$. (4) $\geq 4h$ indicates that the computation time is more than 4 h and the optimal solution cannot be found. (5) (*) indicates that the value is the result of the experimental calculation of 4 h. (6) (-) represents that the value cannot be calculated.

530 discount can greatly improve the profits of the entire supply chain network. As manufacturers
531 receive incremental quantity discount when purchasing raw materials, there is an impact on the
532 purchase volume per purchase. On the small scale problems, it can increase profits by about 3%-
533 11%. On the medium scale problems, the profit can be increased by about 19%-33%. On the large
534 scale problems, it can increase profits by about 31%-38%. Therefore, for logistics network activities
535 with a long planning period, price discounts provided by suppliers can greatly improve the profits
536 of the entire supply chain network.

537 In addition to the impact of discount level on network profits in Table 11, we can also see the
538 solution efficiency of the two algorithms at different scales. The outer-approximation method uses
539 piecewise linear functions to approximate the calculation, and the solution obtained is not accurate.
540 However, because smaller error values are chosen, for small-scale instances, the difference between
541 the solution obtained using the outer-approximation method and the exact solution derived from
542 SOCP is minimal. In small and medium-sized experiments, the solution error value under the two
543 schemes using the outer-approximation method and the SOCP is within 1%. However, for large-
544 scale instances, SOCP cannot provide a solution within a reasonable time. For instances 17-24,
545 SOCP results are limited by a 4-hour compute time limit. In this case, the error value of the two
546 algorithms is controlled within 7% under the two schemes. However, due to the long solution time of
547 SOCP, the outer-approximation method is preferred to solve large-scale problem instances. Overall,
548 as the size of the comparison instances increase, the differences between the two algorithms become
549 more and more significant. Due to the NP-hard of problem (see, e.g., Ahmadi-Javid & Hoseinpour,
550 2015b; Fattahi et al., 2015), the computational efficiency of the outer-approximation method is more
551 advantageous for larger instances, especially when SOCP is difficult to provide a solution within a
552 reasonable time.

Table 11

Comparison of experimental results of two mathematical models for different scales.

Scales	Instances	Discount level		Non-Discount level		GAP _{OA} (%)	GAP _{SOCP} (%)
		OBJ_{OA-1}	OBJ_{SOCP-1}	OBJ_{OA-2}	OBJ_{SOCP-2}		
Small scales	1	123174	123110	112475	112248	9.512	9.677
	2	191598	191432	181536	180897	5.543	5.824
	3	198030	197873	192591	192353	2.824	2.87
	4	265531	265381	252403	251959	5.201	5.327
	5	275589	275448	247303	247040	11.438	11.499
	6	345659	345498	324215	323884	6.614	6.673
	7	377950	377815	357686	357149	5.665	5.786
	8	437017	436757	413947	413488	5.573	5.627
Middle scales	9	234684	233512	185105	184433	26.784	26.611
	10	336636	336352	265182	264502	26.945	27.164
	11	350508	350206	263976	263300	32.78	33.006
	12	487126	486835	393326	392635	23.848	23.992
	13	489126	488848	377904	377219	29.431	29.593
	14	569665	569189	461111	460631	23.542	23.567
	15	589043	587256	474663	473702	24.097	23.972
	16	674985	674122	568995	568252	18.628	18.631
Large scales	17	315875	298542	238495	229156	32.445	30.279
	18	594788	575853	444029	435056	33.953	32.363
	19	680302	659475	503654	499863	35.073	31.931
	20	751012	721217	557763	549765	34.647	31.186
	21	708496	669917	519251	513265	36.446	30.521
	22	852931	819432	649691	642651	31.283	27.508
	23	899782	850263	664633	656210	35.38	29.572
	24	1102087	1089521	798300	790280	38.054	37.865

Notes: (1) OBJ_{OA-1} represents the objective values for our model considering the discount level using outer-approximation algorithm. (2) OBJ_{SOCP-1} represents the objective values for our model considering the discount level using SOCP method. (3) OBJ_{OA-2} represents the objective values for our model without considering the discount level using outer-approximation algorithm. (4) OBJ_{SOCP-2} represents the objective values for our model without considering the discount level using SOCP method. (5) GAP_{OA} is calculated by: $(OBJ_{OA-1} - OBJ_{OA-2}) / OBJ_{OA-2}$. (6) GAP_{SOCP} is calculated by: $(OBJ_{SOCP-1} - OBJ_{SOCP-2}) / OBJ_{SOCP-2}$.

553 From the perspective of model results and algorithm performance, in highly competitive envi-
554 ronments, considering price-sensitive demand contributes to a better understanding of the market
555 and consumers by enterprises, enabling them to formulate more suitable pricing and marketing
556 strategies. This, in turn, increases sales revenue and profits, allowing them to gain a sustainable
557 competitive advantage in fiercely competitive markets. This underscores the necessity of our study
558 in treating price and demand as decision variables.

559 In the context of research on price-sensitive demand, Fattahi et al. (2015) and Ahmadi-Javid
560 & Hoseinpour (2015a) have done some research. However, they encountered limitations in directly
561 obtaining decision solutions for both price and demand variables. On the other hand, while Mo-
562 gale et al. (2022) also treated demand and price as decision variables, their pricing scheme was
563 not a multi-period dynamic decision. Building upon this, we further consider incremental quantity
564 discount at the upstream manufacturer’s procurement stage. Based on the results, it is evident
565 that considering incremental quantity discount creates opportunities for profit growth in the supply
566 chain network. In terms of algorithmic solutions, the SOCP and outer-approximation method have
567 contributed to filling the research gap in handling such NP-hard problems.

568 **6. Conclusion and Outlook**

569 This paper studies a three-echelon multi-period SCND problem that considers both price and de-
570 mand. At the same time, the impact of incremental quantity discount offered by the suppliers on the
571 profitability of the supply chain network is considered. This study also assumes prices as decision
572 variables and establishes an MINLP model with the goal to maximize the profit of the supply chain
573 network. Due to the nonlinearity of the objective function, first, the proposed MINLP model is an-
574 alyzed and the conclusion is obtained that the model [MINLP] has an optimal solution. We further
575 get an important property of this optimal solution that the actual amount of purchased product is
576 equal the demand. This allows us to use SOCP to solve the problem directly through commercially
577 available software such as CPLEX and obtain an exact solution. Meanwhile, these conclusions are
578 vital to solve the problem by outer-approximation method. In addition, both the upper and lower
579 bounds for the number of tangent lines to the revenue function are calculated. Finally, numerical
580 experiments have been conducted to validate the significance of the model and the efficacy of the
581 developed algorithms.

582 The most important findings and management insights are the following:

- 583 (1) The research results provide direct decision-making solutions for product pricing and customer
584 demand. This enables managers to make more accurate pricing decisions by responding to
585 changes in demand while determining location and inventory decisions.
- 586 (2) Considering incremental quantity discount can greatly improve the profit of the whole supply
587 chain network. Especially for logistics network activities with a long planning period, the profit

588 improvement is more obvious.

589 (3) The SOCP algorithm can find the optimal solutions for small-scale and medium-scale problems.
590 However, it cannot solve large-scale problems.

591 (4) The solution of the outer-approximation algorithm is close to the optimal solution for small-scale
592 and medium-scale problems. Moreover, it can still obtain approximate optimal solutions within
593 a reasonable computing time for large-scale problems.

594 Furthermore, there were certain limitations in conducting this study, which can serve as direc-
595 tions for future research efforts. As a primary suggestion, uncertainty methods, such as stochastic
596 optimal control (Gökgöz & Öktem, 2021; Savku & Weber, 2022), robust optimization (Özmen et al.,
597 2017; Khalilpourazari & Hashemi Doulabi, 2022; Goli et al., 2023; Lotfi et al., 2023; Tirkolaee et al.,
598 2023), and grey systems (Gökgöz & Öktem, 2021; Masoomi et al., 2022), can be applied to models
599 to improve the robustness. From the perspective of computation, the effectiveness of linear approx-
600 imation (Taylan et al., 2018) and the more effective accurate algorithms can be further considered
601 to provide the optimal solutions to large-scale instances within a reasonable time.

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605 **Data Availability Statemen**

606 The data that support the findings of this study are available from the authors, upon reasonable
607 request.

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